

Locating-Chromatic Number of Amalgamation of Stars

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Abstract. Let G be a connected graph and c a proper coloring of G . For $i=1,2,...,k$ define the color class C_i as the set of vertices receiving color i . The color code $c_{\Pi}(v)$ of a vertex v in G is the ordered k -tuple $(d(v, C_1), ..., d(v, C_k))$ where $d(v, C_i)$ is the distance of v to C_i . If all distinct vertices of G have distinct color codes, then c is called a locating-coloring of *G*. The locating-chromatic number of graph *G*, denoted by $\chi_L(G)$ is the smallest k such that G has a locating coloring with k colors. In this paper we discuss the locating-chromatic number of amalgamation of stars $S_{k,m}$. $S_{k,m}$ is obtained from k copies of star $K_{1,m}$ by identifying a leaf from each star. We

also determine a sufficient condition for a connected subgraph $H \subseteq S_{k,m}$ satisfying $\chi_L(H) \leq \chi_L(S_{k,m})$.

Keywords: *amalgamation of stars; color code; locating-chromatic number.*

1 Introduction

Let G be a finite, simple, and connected graph. Let c be a proper coloring of a connected graph G using the colors $1, 2, \ldots, k$ for some positive integer k , where $c(u) \neq c(v)$ for adjacent vertices u and v in G. Thus, the coloring c can be considered as a partition \prod of $V(G)$ into color classes (independent sets) C_1, C_2, \ldots, C_k , where the vertices of C_i are colored by *i* for $1 \le i \le k$. The *color code* $c_{\Pi}(v)$ of a vertex v in G is the ordered k -tuple $(d(v, C_1), ..., d(v, C_k))$ where $d(v, C_i) = \min\{d(v, x) | x \in C_i\}$ for $1 \le i \le k$. If all distinct vertices of G have distinct color codes, then *c* is called a *locating-coloring* of *G* . A *minimum locating-coloring* uses a minimum number of colors and this number is called the *locating-chromatic number* of graph G , denoted by $\chi_L(G)$.

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The following definition of an amalgamation of graphs is taken from [3]. For $i = 1, 2, \dots, k$, let G_i be a graph with a fixed vertex v_{oi} . The *amalgamation* Amal (G_i, v_{oi}) is a graph formed by taking all the G_i 's and identifying their fixed vertices. In this paper, we consider the amalgamation of stars. More precisely, for $i = 1, 2, ..., k$, let $G_i = K_{1,n_i}$, $n_i \ge 1$ where v_{oi} be any leaf (a vertex of

degree 1) of K_{1,n_i} . We denote the amalgamation of k stars K_{1,n_i} by

 $S_{k,(n_1,n_2,...,n_k)}$, $k \ge 2$. We call the identified vertex as the *center* (denoted by *x*), the vertices of distance 1 from the center as the *intermediate vertices* (denoted by l_i ; $i = 1, 2, \dots, k$), and the *j*-th leaf of the intermediate vertex l_i by l_{ij} ($j = 1, 2, \dots, m-1$). In particular, when $n_i = m$, $m \ge 1$ for all i, we denote the amalgamation of k isomorphic stars $K_{1,m}$ by $S_{k,m}$.

The locating-chromatic number was firstly studied by Chartrand *et al*. [1]. They determined the locating-chromatic numbers of some well-known classes such as paths, cycles, complete multipartit graphs and double stars. Apart of paths and double stars, the locating-chromatic number of other trees are also considered by Chartrand *et al.* [2]. They constructed a tree of order $n \ge 5$ with the locating- chromatic number k, where $k \in \{3, 4, ..., n-2, n\}$. They also showed that no tree on n vertices with locating-chromatic number $n-1$.

Based on the previous results, locating-chromatic number of amalgamation of stars have not been studied. Motivated by this, in this paper we determine the locating-chromatic number of amalgamation of stars.

Beside that, we also discuss the monotonicity property of the locatingchromatic number for the class of amalgamation of stars. Clearly, the locatingchromatic number of a star $K_{1,n}$ is $n+1$, for any n (since all vertices must have

different color codes). Since any connected subgraph H of $K_{1,n}$ is also a star

with small size, then we clearly have $\chi_L(H) \leq \chi_L(K_{1,n})$. However in general for any connected subgraph $H \subseteq G$, the locating-chromatic number of H may not be necessarily smaller or equal to the locating-chromatic number of *G* .

In this paper, we also investigate the monotonicity property of the locatingchromatic number for amalgamation of stars, $S_{k,m}$. We derive a sufficient condition for a connected subgraph $H \subseteq S_{k,m}$ satisfying $\chi_L(H) \leq \chi_L(S_{k,m})$.

The following results were proved by Chartrand et al. in [1] . The set of neighbours of a vertex v in G is denoted by $N(v)$.

Theorem 1.1. Let c be a locating-coloring in a connected graph G. If u and *v are distinct vertices of G such that* $d(u, w) = d(v, w)$ *for all* $w \in V(G) - \{u, v\}$, then $c(u) \neq c(v)$. In particular, if *u* and *v* are non adjacent *vertices of G* such that $N(u) = N(v)$, then $c(u) \neq c(v)$.

Corollary 1.1. *If G is a connected graph containing a vertex adjacent to k leaves of* G , then $\chi_L(G) \geq k+1$.

2 Main Results

We first prove some lemmas regarding the properties of locating-chromatic number of amalgamation of stars. From now on $S_{k,m}$ denotes the amalgamation of *k* isomorphic stars $K_{1,m}$.

Lemma 2.1. For $k \geq 2, m \geq 2$, let c be a proper coloring of $S_{k,m}$, using at *least m colors*. *The coloring c is a locating-coloring if and only if* $c(l_i) = c(l_n)$, $i \neq n$ *implies* $\{c(l_{ij}) | j = 1, 2, ..., m-1\}$ *and* $\{c(l_{nj}) | j = 1, 2, ..., m-1\}$ *are distinct*.

Proof. Let $P = \{c(l_{ij}) | j = 1, 2, ..., m-1\}$ and $Q = \{c(l_{nj}) | j = 1, 2, ..., m-1\}$. Let c be a locating-coloring of $S_{k,m}$, $k \ge 2, m \ge 2$ using at least m colors and $c(l_i) = c(l_n)$, for some $i \neq n$. Suppose that $P = Q$. Because $d(l_i, u) = d(l_n, u)$ for $c(l_i) = c(l_n)$, for some $i \neq n$. Suppose that $P = Q$. Because $d(l_i, u) = d(l_n, u)$ for
every $u \in V \setminus \{l_{ij} | j = 1, 2, ..., m-1\} \cup \{l_{nj} | j = 1, 2, ..., m-1\}$ then the color codes of l_i and l_n will be the same. So c is not a locating-coloring, a contradiction. Therefore $P \neq Q$.

Let Π be a partition of $V(G)$ into color classes with $|\Pi| \geq m$. Consider $c(l_i) = c(l_n)$, $i \neq n$. Since $P \neq Q$, there are color x and color y such that $(x \in P, x \notin Q)$ and $(y \in P, y \notin Q)$. We will show that color codes for every $v \in V(S_{k,m})$ is unique.

• Clearly, $c_{\Pi}(l_i) \neq c_{\Pi}(l_n)$ because their color codes differ in the *x*thordinate and *y* th-ordinate.

If $c(l_{ij}) = c(l_{ns})$, for some $l_i \neq l_n$, we will show that $c_{\Pi}(l_{ij}) \neq c_{\Pi}(l_{ns})$. We divide into two cases.

Case 1. If $c(l_i) = c(l_n)$ then by the premise of this theorem, $P \neq Q$. So $c_{\Pi}(l_{ii}) \neq c_{\Pi}(l_{ns})$.

Case 2. Let $c(l_i) = r_1$ and $c(l_n) = r_2$, with $r_1 \neq r_2$. Then $c_{\Pi}(l_{ij}) \neq c_{\Pi}(l_{ns})$ because their color codes are different at least in the r_1 th-ordinate and r_2 th-ordinate.

- If $c(l_i) = c(l_{nj})$, $l_i \neq l_n$, then $c_{\Pi}(l_i)$ contains at least two components of value 1, whereas $c_{\Pi}(l_{nj})$ contains exactly one component of value 1. Thus $c_{\Pi}(l_i) \neq c_{\Pi}(l_{nj})$.
- If $c(x) = c(l_{ij})$, then color code of $c_{\Pi}(x)$ contains at least two components of value 1, whereas $c_{\Pi}(l_{ij})$ contains exactly one component of value 1. Thus $c_{\Pi}(x) \neq c_{\Pi}(l_{ij})$.

From all above cases, we see that the color code for each vertex in $S_{k,m}$ is unique, thus c is a locating-coloring.

Lemma 2.2. Let c be a locating coloring of $S_{k,m}$ using $m+a$ colors and $(a) = (m+a-1) \binom{m+a-1}{m-1}, a \ge 0$ $m + a$ $H(a) = (m + a - 1) \binom{m + a - 1}{m - 1}, a$ $\begin{pmatrix} m+a-1 \\ a \end{pmatrix}$ $=(m+a-1)\binom{m+a-1}{m-1}, a \ge 0.$. *Then* $k \leq H(a)$.

Proof. Let c be a locating-coloring of $S_{k,m}$ using $m + a$ colors. For fixed i, let $c(l_i)$ be a color of intermediate vertex l_i , then the number color combinations can be used by $\{l_{ij} | j = 1, 2, ..., m-1\}$ is $\binom{m+a-1}{1}$ 1 $m + a$ *m* $+a \binom{m+a-1}{m-1}$. Because one color is used for coloring the center x, there are $(m+a-1)$ colors for l_i , for every $i = 1, 2, \dots, k$. By Lemma 2.1, the maximum number of k is $(m+a-1)$ $\binom{m+a-1}{m-1}$ = H(a) $m + a$ $(m + a - 1)$ $\binom{m + a - 1}{m - 1} = H(a)$ $+a-1$ + $a-1$) $\binom{m+a-1}{m-1} = H$. So $k \leq H(a)$.

The main result of this paper concerns about locating-chromatic number of $S_{k,m}$.

Theorem 2.1. *For* $a \ge 0, k \ge 2, m \ge 2$, *let* $H(a) = (m + a - 1) \binom{m + a - 1}{m - 1}$ $m + a$ $H(a) = (m + a - 1) \binom{m + m}{m}$ $+a-1$ $=(m+a-1)\binom{m+a-1}{m-1}.$ T . *Then,*

$$
\chi_L(S_{k,m}) = \begin{cases} m & \text{for} \\ m+a & \text{for} \end{cases} \quad 2 \le k \le H(0), m \ge 3,
$$

$$
H(a-1) < k \le H(a), a \ge 1.
$$

Proof. First, we determine the trivial lower bound. By Corollary 1.1, each vertex l_i is adjacent to $(m-1)$ leaves, for $i = 1, 2, ..., k$. Thus, $\chi_L(S_{k,m}) \ge m$.

Next, we determine the upper bound of $\chi_L(S_{k,m})$ for $2 \le k \le H(0) = m-1$. Let *c* be a coloring of $V(S_{k,m})$ using *m* colors. Without loss of generality, we can assign $c(x) = 1$ and $c(l_i) = i + 1$ for $i = 1, 2, ..., k$. To make sure that the leaves will have distinct color code, we assign $\{l_{ij} | j = 1, 2, ..., m-1\}$ by $\{1, 2,..., m\} \setminus \{i+1\}$ for any *i*. Then, by Lemma 2.1, we have that *c* is a locating-coloring. Thus $\chi_L(S_{k,m}) \leq m$.

Next, we shall improve the lower bound for the case of *k* such that $H(a-1) < k \le H(a), a \ge 1$. Since $k > H(a-1)$ then by Lemma 2.2, $\chi_L(S_{k,m}) \ge m + a$. On the other hand if $k > H(a)$ then by Lemma 2.2, $\chi_L(S_{k,m}) \ge m + a + 1$. Thus $\chi_L(S_{k,m}) \ge m + a$ if $H(a-1) < k \le H(a)$.

Next, we determine the upper bound of $\chi_L(S_{k,m})$ for $H(a-1) < k \le H(a)$, $a \ge 1$. Without loss of generality, let $c(x) = 1$ and color the intermediate vertices l_i by 2,3,..., $m + a$ in such a way that the number of the intermediate vertices receiving the same color t does not exceed $\binom{m+a-1}{1}$ 1 $m + a$ *m* $\left(m+a-1\right)$ $\begin{pmatrix} m & n \\ m-1 \end{pmatrix}$, for any *t*. We are able to do so because $H(a-1) < k \le H(a)$. Therefore, if $c(l_i) = c(l_n)$,
 $i \ne n$ then we can manage $\{c(l_{ij}) | j = 1, 2, ..., m-1\} \ne \{c(l_{nj}) | j = 1, 2, ..., m-1\}$. $i \neq n$ then we can manage By Lemma 2.1, c is a locating-coloring on $S_{k,m}$. So $\chi_L(S_{k,m}) \leq m+a$ for $H(a-1) < k \leq H(a)$.

The following figures show minimum locating-colorings on $S_{4,6}$ and $S_{9,3}$.

Figure 1 A minimum locating-coloring of $S_{4,6}$.

Figure 2 A minimum locating-coloring of $S_{9,3}$.

Next, we discuss the monotonicity property of locating-chromatic number for the amalgamation of stars.

Theorem 2.2 If $2 \le k \le m-1$, then $\chi_L(G) \le \chi_L(S_{k,m})$ for every $G \subseteq S_{k,m}$ *and* $G \neq K_{1,m}$.

Proof. Let c be a minimum locating-coloring of $S_{k,m}$ obtained from Theorem 2.1. Let G be any connected subgraph of $S_{k,m}$. Define a coloring c' on G by preserving colors used in $S_{k,m}$ for the corresponding vertices, namely $c'(v') = c(v)$ if *v* is the corresponding vertex of *v*' in $S_{k,m}$. We show that *c*' is a locating-coloring of *G* .

If there exist l_i , l_n such that $\{c'(l_{ij}) | j = 1, 2, ..., r\} = \{c'(l_{nj}) | j = 1, 2, ..., r\}$, with $1 \le r \le m-1$, then color codes of l_{ij} and l_{nj} for every $j = 1, 2, 3, \dots, m$ is unique because $c'(l_i) \neq c'(l_n)$ for every $l_i \neq l_n$. If $c'(l_i) = c'(l_{nj}) \neq c'(x)$, then the first component of $c'_n(l_i)$ has value 1, whereas for $c'_n(l_{nj})$ it has value 2. So color code of l_i and l_{nj} are different. Next, if $c'(x) = c'(l_{nj})$, $G \neq P_3$ then their color codes are different because $c'(l_i) \neq c'(l_n)$ for every $l_i \neq l_n$. For the case $G = P_3$, $v_i \in V(P_3)$ for each *i* is colored by 1, 2, and 3 respectively. Because the color codes for every $v \in V(G)$ is unique, then c' is a locating-coloring of G. So $\chi_L(G) \leq \chi_L(S_{k,m})$ for every $G \subseteq S_{k,m}$, $G \neq K_{1,m}$.

Let $S_{k,(n_1,n_2,\ldots,n_k)} \subseteq S_{k,m}$. Define $A = \{i \mid n_i = 1\}$. For $k \ge m$, we must restrict subgraphs of $S_{k,m}$ so that satisfy monotonicity property.

Theorem 2.3 *If* $k \ge m$ *and* $|A| \le \chi_L(S_{k,m}) - 1$ *then* $\chi_L(S_{k,(n_1,n_2,...,n_k)}) \le \chi_L(S_{k,m})$.

Proof. Let $k \ge m$ and from Theorem 2.1, we have that $\chi_L(S_{k,m}) = m + a$ for $H(a-1) < K \le H(a), a \ge 1$. Let $G = S_{k,(n_1,n_2,...,n_k)}$ be any subgraph obtained from $S_{k,m}$ with $1 \le n_i \le m$. If $2 \le n_i \le m$ for each *i*, then color vertices of *G* follow the proof of Theorem 2.1. Clearly, the coloring of *G* is a locating-coloring. Otherwise, we have $n_i = 1$ for some i, and so $|A| \ge 1$. If $|A| \le \chi_L(S_{k,m}) - 1$, then the center *x* is given color 1, $l_i \in A$ for each *i* is colored by 2,3,..., $\chi_L(S_{k,m})$, respectively and the colors for the other vertices follow the proof of Theorem 2.1. Observe that the color codes of l_i for each $l_i \in A$ has value 1 in the 1thordinate, 0 in the *i* th-ordinate, and 2 otherwise, these color codes are unique. For the remaining of the vertices, the color codes are also unique as proven in Theorem 2.1. As the result, the coloring of G is a locating-coloring. So $\chi_L(S_{k,(n_1,n_2,\dots,n_k)}) \leq \chi_L(S_{k,m}).$

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